

Introduction

Motivations for this talk : LLMs strengths

- Strong capabilities in code generation
- Effective for semantic code analysis task
- Operate by learning implicit patterns data



Figure – Overall score of DevQualityEval. Data from <https://devqualityeval.com/>.

Section 2

Can LLMs classify floating-point exceptions in small C functions for given inputs?

InterFLOPBench¹ : Floating-Point Benchmark to Test Numerical Software Analysis Tools

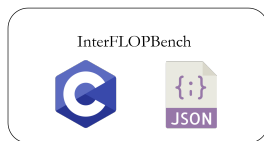


Figure – C source code with accompanying JSON metadata file. A total of **93** benchmarks, **1 150** samples (inputs) and **1 315** errors occurrences.

Category	Occurrences
No Error	478
Comparison	255
Cancellation	214
Overflow	153
Underflow	88
Div. Zero	70
NaN	57

Table – InterFLOPBench Error Category Distribution.

Example of kernel in InterFLOPBench

C file

```
double kernel(double x, double y)
{
    double oscillator = cos(x) - y;
    double logValue = log(oscillator);
    double result = (exp(logValue / 2.0)
        + log(fabs(oscillator) + 1.0))
        / (1.0 + fabs(logValue));

    return result;
}
```

JSON file

```
{
  "ordered_inputs": ["x", "y"],
  "ordered_outputs": ["result"],
  "datasets": [
    {
      "errors": ["no_error"],
      "inputs": [
        {
          "x": 0.0, "y": -0.5,
          "x": 1.57, "y": -0.2
        }
      ],
      "errors": ["nan"],
      "inputs": [
        {
          "x": 0.0, "y": 1.5,
          "x": 3.14, "y": 1.1
        }
      ]
    }
  ]
}
```

Errors
predicted by
FPChecker
for the
inputs below.

File name : 48_complex_log

This example trigger NaN if $\cos(x) \leq y$.

Micro F1 Score for Multilabel Evaluation

7 categories : cancellation, underflow, overflow, NaN, division by zero, comparison, and no error.

Micro F1 Formula

$$F1_{micro} = 2 \times \frac{TP}{2TP + FP + FN}$$

Scenario	True	Predicted	Impact
Correct	overflow	overflow	TP overflow
Missing	overflow	no error	FN overflow + FP no error
False alarm	no error	overflow	FP overflow + FN no error
Under-pred	{overflow, NaN}	{overflow}	TP overflow + FN NaN
Over-pred	{NaN}	{overflow, NaN}	TP NaN + FP overflow

Figure – TP/FP/FN impact. The "True" is obtained from the dynamic exception-based tool FPChecker.

Overall Results

Full per-benchmark disagreement analysis is available online at :
github.com/interflop/InterflopBench/ground_truth_analysis_table.html

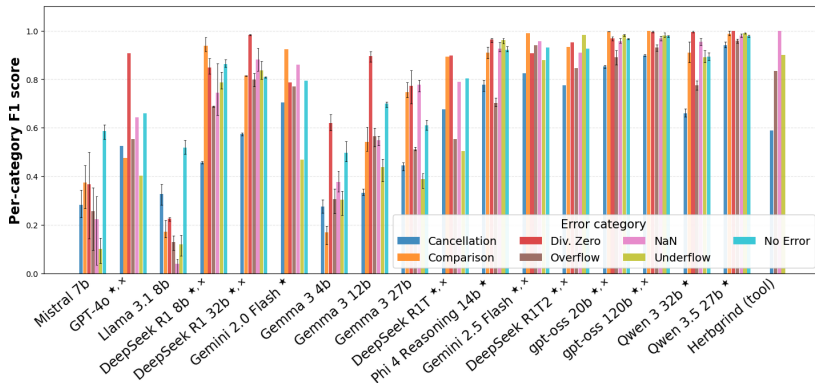


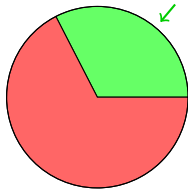
Figure – Per-category F1-score across 17 LLMs sorted by release date from the oldest (left), to the latest (right). * Chain-of-Thought (CoT) reasoning ; x Mixture-of-Experts (MoE) architecture.

Disagreement Analysis & Oracle Refinement

LLM vs. FPChecker : 86 disagreements, 28 LLM victories

- **Instruction isolation :**
FPchecker cancellation checks do not handle FMA operations.
- **Explicit inequality checks :**
FPChecker detects strict equality comparisons (`==`) but overlooks inequality checks (`!=`) between two floating-point variables.
- **Library function propagation :** FPChecker does not instrument `math.h` library functions.

gpt-oss-120b : 28 (32.6%)



FPChecker : 58 (67.4%)

Disagreements
(n=86 out of 1 130 test samples)

Conclusion : LLM already have some basic numerical knowledge

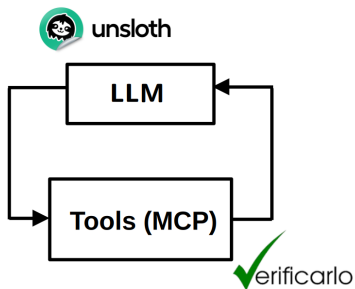
Limitations :

- No formal correctness proof
- Absence of guarantees

Practical applications :

- Identifying likely error patterns
- Prioritizing suspicious code regions for review

MCP Servers with Local LLMs using Unsloth



MCP Verificarlo : compile, instrument, and analyze numerical code

- Browse project files, read source code, and detect Makefiles
- Compile C/C++/Fortran code through Verificarlo compiler
- Run binaries under different floating-point backends
- Compare numerical stability across optimization levels or algorithm variants
- Explore reduced-precision effects
- Detect catastrophic cancellation at runtime
- Delta-debug to detect the exact source lines causing numerical instability

Use Case : Kahan Summation

```
1 REAL sum = f[0];  
2 REAL c = 0.0, y, t;  
3 for (i = 1; i < N; i++) {  
4     y = f[i] - c;  
5     t = sum + y;  
6     c = (t - sum) - y;  
7     sum = t;  
8 }  
9 return sum;
```

Kahan Use Case : Two Reports

Claude Code (VS Code) vs Local model (Qwen 3.5 9b)

	Validation	Local system
IEEE baseline	5.08125470021-33238e+02	5.08125470021-33238e+02
MCA samples	30 runs ; mean 508.1254700213325	30 runs ; “very close to baseline”
Std. dev.	5.55×10^{-13}	<i>not reported</i>
Significant digits	CNH : 14.54 General : 14.45	“typically 15–16” (<i>no method named/ exact value</i>)
Verdict	Stable, no further debugging needed	Stable, no issues detected

Execution Time

Metric	Time	Speed
Prompt eval (s)	130.35	35 tok/s
Generation (s)	401.23	5.4 tok/s
Tokens		2,173
Total wall time	694.91 s (11.6 min)	

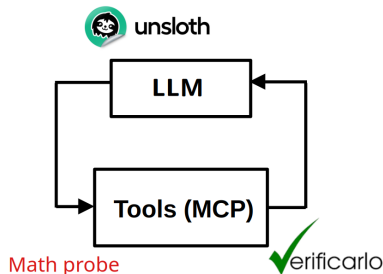
Use Case : Riemann–Siegel

```
1 $ ./RiemannSiegel 10 100000 10
2 I found 137931 Zeros in 6.934 seconds      # INCORRECT
3
4 $ ./RiemannSiegel 10 100000 100
5 I found 138069 Zeros in 56.035 seconds     # OK
```

Expected answer

Counting sign changes of $Z(t)$ on a uniform grid is only correct if the grid is fine enough to resolve *every individual zero*.

MCP Math Probe



- Will this expression overflow the target dtype?
- Is this subtraction losing decimal digits?
- Is this sampling step too coarse for the underlying oscillation?

Qwen 3.5's Conclusion : A Mixed Signal

The final answer contains **two different mechanisms** :

The correct sentence

"The coarser sampling either misses the exact crossing point due to step size = 0.1, evaluates both points on the same side of zero due to undersampling. . . the finer sampling catches the intermediate sign change."

The incorrect sentence, right next to it

"Numerical noise in $Z(t)$ causes sign flips at slightly different t -values. . . within numerical noise ($\sim 10^{-12}$)."

Execution Time

Phase	Time	Speed
Prompt eval (time to first token)	600.28 s (10.0 min)	36.7 tok/s
Generation	2185.51 s (36.4 min)	4.4 tok/s
Tokens generated		9,544
Total wall time	4127.14 s (68.8 min)	

Conclusion & Future Directions

- The fix that matters most for efficient computation : localize before instrumenting
- For scalability, a major open question is how an LLM reads and navigates large source files
- Tools that can localize floating-point-relevant computation in code before deep instrumentation is applied

